

Elementary ARM for Reversing

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Why learning?

- Most of mobile platforms are based on ARM.
 - Android
 - Symbian
 - Windows Phone 7
 - Apple iOS
- As well as some network devices.
- **New trends in malware's evolution**

What's ARM

- A company
- CPU products
- Instruction architectures

Versions

Products	Instruction sets
ARM7	ARM v4
ARM9	ARM v5
ARM11	ARM v6
Cortex-A	ARM v7-A
Cortex-R	ARM v7-R
Cortex-M	ARM v7-M

Architecture

- 32 bits
- RISC
- von Neumann architecture
 - Mixed data and code
 - ARM7
- Harvard architecture
 - Separated data and code
 - ARM9 and later

Registers

- **R0 – R3 (a1 – a4)**
 - Common registers
 - For argument passing and result return
- **R4 – R11 (v1 – v8)**
 - Common registers
 - For local variable
- **R12 (IP)**
 - Intra-Procedure-call scratch register

Registers (cont.)

- R13 (**SP**)
 - Stack Pointer
- R14 (**LR**)
 - Link Register
- R15 (**PC**)
 - Program Counter

Registers (cont.)

- PSRs: Program Status Registers
 - ALU flags (zero flag and carry flag)
 - execution status
 - current executing interrupt number

Data instructions

MOV*cdS* reg, arg

CMP*cd* reg, arg

ADD*cdS* reg, reg, arg

Form:

<opcode> { <cond> } { S } <Rd>, <Rn> { , <operand2> }

Data instructions (cont.)

- `opcode`: `MOV`, `ADD`, `CMP` ...
- `cond`: `EQ`, `NE`, `GE`, `LT` ...
- `s`: update PSR flags
- `reg`: `R0` – `R15`
- `arg`: right-hand argument (next page)

Forget ALL above!

Just lookup references when using them.

Right-hand argument

- `#imm8`: immediate (rotated into 8 bits)
 - `MOV R0, #12`
- `reg`: register
 - `ADDS R0, R0, R1`
- `reg, shift`: register shifted by distance
 - `CMP R0, R1, LSL #2`
 - `ADDNE R0, R0, R1, ASR R3`

Branch instructions

- **B***cd* imm₁₂
 - Branch to imm₁₂ words away
- **BL***cd* imm₁₂
 - Copy PC to LR, then branch
- **BX***cd* reg
 - Copy reg to PC

Memory instructions

- **LDR***cd***B** *reg, mem*
 - Loads word/bytes from memory
- **STR***cd***B** *reg, mem*
 - Stores word/bytes to memory
- **B**: operate by **b**yte instead of by word
- *mem*: memory address (next page)

Memory address

- $[reg, \# \pm imm_{12}]!$
 - reg offset by constant
- $[reg, \pm reg]!$
 - reg offset by variable bytes
- $[reg_a, \pm reg_b, shift]!$
 - reg_a offset by shifted variable reg_b
- $!:$ update reg , then access memory

Memory address(cont.)

- $[reg], \# \pm imm_{12}$
 - Access address reg , then update reg by offset
- $[reg], \pm reg$
 - Access address reg , then update reg by variable
- $[reg], \pm reg, shift$
 - Access address reg , then update reg by shifted variable

Memory instructions(cont.)

- **LDM***cdum* *reg!* , *mreg*
 - Loads memory in *reg* into multiple registers
- **STM***cdum* *reg!* , *mreg*
 - Stores multiple registers into memory in *reg*
- *um*: update mode (next page)
- **!**: change the register
- *mreg*: e.g. {**R0-R3** , **R7** , **R10**}

Update mode

- **IA**: address **I**ncrement **A**fter each read
- **DA**: ... **D**ecrement **A**fter ...
- **IB**: ... **I**ncrement **B**efore ...
- **DB**: ... **D**ecrement **B**efore ...

Stack model

- Full Descending stack
- Push: `STMFD SP!, mreg ; ==STMDB`
 - `SUB SP, SP, 4`
 - `MOV Rn, [SP]`
- Pop: `LDMFD SP!, mreg ; ==LDMIA`
 - `MOV Rn, [SP]`
 - `ADD SP, SP, 4`

More ATPCS

- Parameter passing
 - Use **R0 – R3 (a1 – a4)** firstly, copy from left to right;
 - Use **stack** if more than 4, push from right to left.
- Result return
 - Use **R0 – R3 (a1 – a4)** firstly
 - Use an additional **address parameter** if 4 registers is not enough

Example

```
51 void restart_adb(pid_t pid)
52 {
53     kill(pid, 9);
54 }
```

Example(cont.)

restart_adb

; CODE

pid = -8

STMFD SP!, {R11,LR}

ADD R11, SP, #4

SUB SP, SP, #8

STR R0, [R11,#pid]

LDR R0, [R11,#pid] ; pid

MOV R1, #9 ; sig

BL kill

SUB SP, R11, #4

LDMFD SP!, {R11,PC}

; End of function restart_adb

```
return 0;
```

Thank you!

Q & A